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FROM THE EDITOR

Please contact me if you were at the 'RT-11 Can Do It' session at San Diego and have the session taped. I would like to include parts of the session in the next newsletter.

Thankyou,

Ken

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USER INPUT

Chester & Lois Wilson
Phone (0745911)-201

71 Galatea St,
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Australia.
17-Oct-80

Dear Ken,

I have had some trouble getting much interest in Runoff from your side of the Pacific, but the efforts of Bob Denny (if I read his signature correctly) in the last Mini-Tasker have prompted me to ask if you could please add this letter to your publication.

The version of RUNOFF (NO1) which I submitted to DECUS some months ago has prompted a wish list (see below), and a couple of fairly easily correctable bugs have been found by David O'Connor (of Buffalo). Here they are.

PROBLEM 1: No Hyphenation

This is a source-level patch in HYPHEN.MAC. Calls to ALPH and ALPH2 are tested on return with BEQ and BNE instructions, instead of BCS and BCC instructions. Here is the TECO command string for correction of this, with dollar signs being escapes, and "A" being control-A.

```
.R TECO
*ebHYPHEN.MAC$$
*<nPC,ALPH$;
lcc
::fsNE$CS$*u::fsEQ$CC$*u
"A? Error in HYPHEN.MAC
"A2v'>$$
*ex$$
```

PROBLEM 2: /P repeats Headers

This is an error of omission on my part. The routine in RNRT11.MAC for getting a line from the user's terminal (TTC4N) is missing a RETURN statement at the end, with the result that execution falls straight through into TTC3, sending the header (if one is present) to the output file before it actually should get there.

The TECO instructions for correcting this are as follows:

```
.R TECO
*ebRNRT11.MAC$$
*nTTC4N:::n.PAGE$OL$;
*<tab>RETURN
$$
*ex$$
```

(The <tab> is a single tab character.)

".special <delimiter><characters><delimiter>" with control characters being permissible within this string. It may be useful particularly for screen use to have strings predefined (prefix file) for different terminals and called merely by ".special <name>" where <name> has been defined in a ".spcdef /string/" as per .special.

6. It would be useful to have a ".tell <text>" to type information to the user during processing. Similarly a ".wait" - for example, to allow changing daisy wheels or golf balls on a printer.
7. Some ability to cope with proportional spacing printers, for example by reading in a table of divisions for each character.
8. An ability to underline blank spaces.
9. Ability to put page numbers on either (or alternating) sides of the page, or in the middle at the bottom.
10. (This is a crazy one to see what you think): A relatively simple form of conditional processing could be implemented now that it is possible to have multiple file input. Consider:

I am aiming to collect assorted ideas and bugs, and re-release to DECUS in about 6 to 12 months. The original aim had been to leave the source code completely alone so as to maintain RSX compatibility, but after the HYPHEN source bug was detected I reckon extensions are definitely in order, keeping the sources so that they may be passed across to anyone from RSTS/RSX who wants to upgrade!

Current gripes/bugs: (to solicit those of EVERYONE interested please!)

1. Assorted blank pages at end of output (probably will tackle by holding output in RNRT11 and ignoring blank pages there)
2. .noheader leaves a blank line at the top of a page.
3. .autoparagraph causes a bombout if there are blank lines at the end of the input file. It also causes a break performs an implicit .paragraph, which is most undesirable if it is to be included in a prefix .RNO file. For compatibility it may be necessary to have a ".setap" command, which just sets up the parameters as per .ap but doesn't have the additional effect.
4. .literal fails to include blank lines. I feel this is a damned nuisance.
5. A mechanism should be included to allow addition of special characters (eg for setting up LA120 or Diablo printers), with no recording of them from the point of view of horizontal or vertical counting. It could take the format of:

```
.SETTRUE PARAM1 (in prefix file)
...
.IFTRUE PARAM1
<runoff commands etc>
.IFEND (?ENDC PARAM)
(similarly .IFFALSE)
```

I would appreciate any input you have re all this. Runoff is getting a fair pounding this side of the Pacific, and the better we can make it the more use to everyone.

Any more bugs anyone can find will be gratefully appreciated.

Regards,

Chester Wilson
[Chester Wilson]



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DEPARTMENT OF RADIOLOGY
Radiation Physics and Safety Office

Wayne A. Meyers, M.S., Director
Michael G. Stambaugh, M.S.
John G. Wolodzko, B.S.

6 November, 1980

Ken Demers
MS-48
United Technologies Research Centre
East Hartford, Connecticut 06108

Dear Mr Demers,

The following note is intended to spare some grief for anyone who may initialize a disc, by design or otherwise, only to find that their backup volume is unusable.

As you may be aware, section 8.2.12.5 of the System User's Guide for RT11 version 4 describes the DUP option switch /D to "uninitialize" a volume immediately following an initialization. But, on the occasions I've had to employ this technique, it has been completely ineffectual, the only result being the error message: "DUP-F-INVALID RESTORE DATA-DK1:". This means that the device Initialize/Restore Data Area at locations 204-251 of the home block does not contain a duplicate copy of the device directory (for an unexplained reason) which existed before the initialization.

Now, contrary to the statement accompanying the error message, all files and directory entries existing prior to the initialization are not lost. Initialization does not delete the original device directory. The directory segment link words are altered, however, in addition to other changes, to indicate that no segments are in use. By employing the save image patch utility, SIPP, described in chapter 22 of the RT-11 version 4 System Reference Manual, it is possible to patch the device directory segment link words, as described in section 9.1.6.2 of the Software Support Manual. A careful review of chapter 9 makes the process quite clear and simple.

A final note of caution: Apparently the first few files on a disc can be lost during an initialization, but most of the user's files, if not all of them, will be listed in the directory subsequent to the above procedure. Other parameters, such as the total number of blocks, or free blocks, on the device will be incorrect, but these factors do not interfere with the transfer of files to another disc.

Yours sincerely,



John G. Wolodko



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October 7, 1980

Mr. Ken Demers,
MS-48,
United Technologies Research
Center,
East Hartford, CT
USA 06108.

Dear Mr. Demers:

Some years ago I submitted descriptions to the 'Mini-Tasker' and the 'Multi-Tasker' of two RT-11 related systems which have since found popular acceptance in the community. These were:

RDCL: A system which allows an RT-11 computer to transfer files to and from an RSX-11M system and also provides 'virtual terminal' support from the RT system to the 11M system;

RTV: An implementation of RT-11 V02B which runs as a task under RSX-11M.

Development of the 'RDCL' system has continued since that time. To date, there are over 50 installations worldwide which run RDCL. Some installations use RDCL to support a large number of 'core-only' LSI-11's running RT-11 for production, dedicated-type applications. Others use RDCL as a means of augmenting the mass storage of their disk-based RT-11

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machines by using 'virtual disk' space on the 11M machine. Most of the users have reported that the system runs very well, is efficient, and constitutes a viable alternative to DECNET. Users who would like more information on the latest release can receive a full documentation set (free of charge) by writing to me.

RTV was originally the brainchild of Mr. Bill Clogher, then a DEC systems programmer, who, in his own words "developed a prototype one evening after a party" in order to convince DEC management that such a product was easily implemented and that a market would exist for it. Bill handed over the system to me at a DECUS meeting in 1976 after it became evident that management was not interested. I did a bit of development on it, and was using it in-house for a while in addition to distributing it to interested users.

Because the original design of RTV was 'quick and dirty', it

proved to be an impossible job to upgrade it when RT-11 was upgraded from V02B to V03. However, user interest was still (and is even now) high and so I continued to distribute it. Several unsuccessful attempts were made at U.S. DECUS meetings to convince DEC that a market does exist for the product, but to little avail. We got vague hints and 'don't quote me' promises that funding for it existed and that it was just a matter of time before it became available. I have finally given up.

My purpose in writing this is twofold:

- a) I wish to inform interested users that I will no longer distribute RTV; and
- b) I wish to encourage interested users to send purchase orders to DEC for "RTV V01.0" as a method of demonstrating that there is a market for such a product. I would suggest that they make purchases conditional upon 90 days delivery.

Judging by the number of enquiries which I receive for 'RTV', I'm fairly convinced that this strategy may achieve a useful purpose. It would be interesting to get an 'official' DEC reply to this letter published in some future issue of the Newsletter.

Thank you for the space.

Yours truly,



Gary Bernstein,
Manager, Computer Systems,
Biomedical Engineering Unit.

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PROBABILITY	MEAN OF D	STANDARD ERROR
-------------	-----------	----------------

0.1	0.000074	0.000004
0.05	0.000134	0.000010
0.01	0.000202	0.000013
0.001	0.004739	0.003059
0.0001	0.062081	0.040453
0.00001	0.072889	0.044449

LISTING OF THE PROGRAMS : (SUBROUTINES PROBI AND NDTRI).

RTRAN IV V01C-03A

001 SUBROUTINE PROBI(F,P,N1,N2,IER)

APPROXIMATION OF THE VALUE OF THE VARIANCE RATIO F ASSOCIATED WITH
A PROBABILITY P AND N1 AND N2 DEGREES OF FREEDOM.

002 REAL*8 A,B,C,D,AA,BB,CC,DELTA,X

CALCULATION OF THE CRITICAL VALUE Z OF THE NORMAL DISTRIBUTION.

003 CALL NDTRI(P,Z,D,IER)

004 C IF(Z.LT.0)Z=-Z

THIS APPROXIMATION IS VALID FOR N2 VALUES GREATER THAN OR EQUAL
TO THREE.

006 C IER=0

007 C D=2./(<9.*N1)

008 C=2./(<9.*N2)

009 A=1-C

010 R=1-D

011 C AA=C*Z-Z-A*A

012 BB=2*A*B

013 CC=-B*B+D*Z*Z

014 C DELTA=BB**2-4*AA*CC

015 C IF(DELTA>1,2,3

016 1 IER=1

017 RETURN

018 2 X=-BB/(2*AA)

019 GO TO 4

020 3 X=(-BB-DSQRT(DELTA))/(2*AA)

021 4 IF(X.GT.0)GO TO 5

023 IER=2

024 RETURN

025 5 F=X*X*X

026 RETURN

027 END

RTRAN IV V01C-03A

001 SUBROUTINE NDTRI(P,X,D,IER)

DETERMINATION OF THE Z VALUE (X) AT A GIVEN LEVEL OF PROBABILITY P.

REFERENCES : ABRAMOVITZ M. AND STEGUN L.A., HASTINGS C. (1955)

002 A0=2.515517

003 A1=0.802853

004 A2=0.010328

005 B1=1.432788

006 B2=0.189269

007 B3=0.001308

008 PP=0.3989423

009 C

010 IER=0

011 IF(P>1,4,2

012 1 IER=-1

013 GO TO 12

014 2 IF(P-1.)>7,6,1

015 4 X=-1.E25

016 5 D=0.

017 GO TO 12

018 6 X=+1.E25

019 GO TO 5

020 7 D=P

021 IF(D-0.5)>9,9,8

022 8 D=1.-D

023 9 T2=ALOG(1./(<D*D))

024 T=T2**0.5

025 X=T-(A0+A1*T+A2*T2)/(1.+B1*T+B2*T2+B3*T*T2)

026 10 IF(P-0.5)>10,10,11

027 11 X=-X

028 D=PP*EXP(-X*X/2.)

029 12 RETURN

030 END

The followings were developed by Alan Mazer and Ron Fires.
Questions about them can be answered by sending a SASE to:

Alan Mazer
U.S. Borax Research Corporation
412 Crescent Way
Anaheim, CA 92802

SQUEEZING WITHOUT DUP

=====

There is now a way to do a type of device squeeze without
really squeezing. The reason for developing this can be seen by
looking at the way we use RT-11. We use the FB monitor to run a
foreground job that reads data from three automated instruments, and
copies the data from each to their associated TTY and associated
data file on a floppy. This means the foreground has three files

open at one time, and, since each instrument is at a different stage of analysis, stopping all three instruments to allow a squeeze would be time consuming. The solution is to copy files from the end of the disk to free space in the middle. Usually, moving a few files is enough to make a large enough free space. Back in RT-11 V2, PIP had an X option that would copy a file to itself, only at a different place on the disk. This was how a file could be moved to free up space. As PIP was "enhanced", the X option was discarded. We have brought back a way of moving files which is implemented in TECO.

The first part of the squeeze package is the file MOVE.TEC. This TECO macro takes as an argument the name of the file to be moved. The macro reads the file into the buffer, sizes it, then allocates a file that is the correct size, and finally writes out to the file. The program uses the same file name, and does not make a backup. The means that a warning message (%Superseding existing file) is printed, but since there is no backup to delete, your space is already freed. This macro is limited because it only works on ASCII text files (no SAV or OBJ files) and the file must be small enough to fit in the buffer (this is so MOVE can determine its size). The macro will print an error message if the file is too big, or the file can not be found. You can use this macro to move one file at a time. It is used as follows:

Lets say you need 30 blocks, and your disk is fragmented. Here is part of a full directory listing :

```
< UNUSED >      14
WASTE .SPC       7
< UNUSED >      20
```

If you save the monitor command:

```
.ED/TEC/EXE:WASTE.SPC MOVE
```

You would then have the following directory:

```
WASTE .SPC       7
< UNUSED >      34
```

You now have more than 30 blocks. As you can see, the macro is a useful tool for quick moves, but if you needed to do several, it might take a while. A way of automating it further is needed.

The next step is to do MOVES on enough files to get the disk a little less fragmented. We found that, in general, the last part of the directory is the most fragmented. That means that a move of the last few files would clear up enough space, and this is, in our case, what happens. To automate this process, there is a TECO program, called SQUEEZ.TEC, that will read a directory listing from a file, pick up the number of files you want moved from the end of the directory, and generate a .COM file to execute the MOVES on those files. When the program exits, it also starts the new command file. The number of files to be moved is determined by either by: 1) If the ED/EXE command has an argument (e.g. ED/EXE:10), that is used, or 2) If there is no argument, the program will ask the user how many files to move.

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SQUEEZ changes a directory listing, in FILE.DIR, into the command file. It assumes that the listing has been made with the options /FAST and /COLUMN:1. To save time and thinking, a SQUEEZ.COM file was written that contains the directory command, the SQUEEZ macro start-up, and some other commands. We simply enter @SQUEEZ and its off and running. With these three files, we can do partial squeezes without interrupting the foreground.

HOW THEY WORK (briefly)
=====

MOVE: Strip block size ([n]) from file name, if given.
Put file name in F.
Make file name into macro to read file, and put in O.
Kill buffer. Execute O (open file). Read into buffer.
Save pointer to end in I. Append.
Compare old pointer to new end, if different, file was too big.
Determine file size. Make write macro that allocates correct size.
Execute write macro. Exit.

SQUEEZ: Initialize counter C (number of files to move).
If argument was passed, put in C, and tell # of files moved.
If no argument passed, read from TTY and put in C.
Open FILE.DIR for input and output (so no .BAK is created).
Delete bottom two lines (# files and # free blocks in listing).
Move up C lines. Kill from beginning to pointer.
Remove all blanks.
Make each line into a command to ED/EXE:file MOVE.
Insert commands to delete FILE.DIR, give a directory of free space, and set the TT to quiet.
Exit and start FILE.DIR, the command file just created.

```
.R TECO
*ERMOVE.TEC%Y$HT$$
!      EDIT 2 of  on 31-Jul-80 by VT52 TECO !
!      EDIT 1 of  on 31-Jul-80 by VT52 TECO !
!      EDIT 0 of MOVE.TEC[3] on 31-Jul-80 by VT52 TECO !
! CALL SEQUENCE:
```

```
EDIT/EXE:file.ext MOVE
```

```
!
```

```
OJ:SC$'SRK' HXF OJ:ER$ZJ@I/$'F ^A
?MOVE-F-Output file not found - ^A:GF^A
^AHKEKEX'$$/ HXO HK MO Y ZU1 A ZJ Q1-.'N ^A
?MOVE-F-File too big for move
^AHKEX' Q1/512+1U1 OJ1
$OJGFII$Q1\@I/]$$/ OJIEW$OJ 1XC0J1KMC
EX*$C$
```

```
.R TECO
*ERSQUEEZ.TEC%Y$HT$$
!      EDIT 12 of SQUEEZ.TEC on 27-Aug-80 by VT52 TECO !
!      EDIT 0 of SQUEEZ.TEC[2] on 22-Aug-80 by VT52 TECO !
OUC Z'NOJ\UC'
QC'E @O!READIN! '
```

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```

^ANumber of files to be moved = ^A QC=^A
^A @!OPEN!
!READIN! ^ANumber of files to be moved ? ^A
1*2*8*ETET < ^TUD QD-13^E 0^/
QD-127^E QC/10UC 8^T 32^T8^T^
QD^D QD^T (QC*10)+(QD-^0)UC ^ > ^TUD ^A
^A
!OPEN! ERFILE.DIR$Y$EWFILE.DIRC2J$
ZJ-2K
ZJ-QCLB, .K
<:S $;FR$>
QC<IED/EXE:$LRRI MOVE$L>
!DELETE FILE.DIR/NOQUERY
DIR/FREE
SET TT QUIET
$EG@FILE.DIR$$$^C$>

```

```

.TY SQUEEZ.COM
!      EDIT 2 of SQUEEZ.COM on 27-Aug-80 by VT52 TECO !
!      EDIT 1 of  on 22-Aug-80 by VT52 TECO !
!      EDIT 0 of SQUEEZ.COMC2J on 22-Aug-80 by VT52 TECO !
SET TT NOQUIET
DEL *.BAK/L
DIR/FA/C:1/OUT:FILE.DIR
ED/EXE SQUEEZ
DIR/FR

```

.@SQUEEZ

.DEL *.BAK/L
?PIP-F-File not found DK!*.BAK

.DIR/FA/C:1/OUT:FILE.DIR

.ED/EXE SQUEEZ

Number of files to be moved ? 3
%Superseding existing file

.ED/EXE:TS1187,JF2 MOVE

%Superseding existing file

.ED/EXE:TS1187,KFE MOVE

%Superseding existing file

.ED/EXE:SQUEEZ.DOC MOVE

%Superseding existing file

.DELETE FILE.DIR/NOQUERY

.DIR/FREE

16-Sep-80

< UNUSED >	1	< UNUSED >	2
< UNUSED >	2	< UNUSED >	1
< UNUSED >	2	< UNUSED >	2
< UNUSED >	1	< UNUSED >	1

< UNUSED >	3	< UNUSED >	5
< UNUSED >	3	< UNUSED >	2
< UNUSED >	2	< UNUSED >	4
< UNUSED >	5	< UNUSED >	3
< UNUSED >	2	< UNUSED >	2
< UNUSED >	10	< UNUSED >	16
< UNUSED >	0	< UNUSED >	4
< UNUSED >	16	< UNUSED >	26
0 Files, 0 Blocks			
115 Free blocks			

.SET TT QUIET

Ian Hammond
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West Germany

DECUS Symposium, Amsterdam

The Amsterdam Symposium was well attended. The RT-11 training session, given by Anton Chernoff, overflowed its planned seating capacity. The session included tutorials on MTT and XM and some hours of questions and 'real' answers on V4.

The 'What to do when RT-11 runs out' session, planned to look at growth paths for to RT-11, turned into an RT-11/RXS-11 session since only the RT-11 and RXS-11 representatives attended. Some would have liked to hear more about RSTS - particularly since RSTS already has an RT-11 emulator. The official Digital reaction is that the 'RSTS RT-11 emulator is unsupported'. But then, some of Digital's best software is 'unsupported' (e.g. TECO). Indeed, the RT-11 development group use an 11/70 RSTS-E system to develop much of RT-11.

A later session handled RT-11/RXS-11 migration specifically. Ruediger Meisenburg (Heinrich Hertz Institute, Berlin) presented a very clear paper on moving a telephone network simulation from RT-11 to RXS-11. The main issue was mapping RT-11 SYSLIB calls into equivalent RXS-11 calls. This session concluded with a presentation of the RXS-11 RT-11 emulator. Les Parent (RT-11) and Stephen Paavola (RXS-11) stressed (repeatedly) that there was no commitment to actually release the emulator. Digital polled potential support for the emulator - even about pricing. Educated opinion has it that the emulator passed its test of fire and probably will be released. The next version of RXS-11M should finish field-tests and be out in the third quarter of '81 (RXS does a year of field tests). This RXS-11M version will probably support DCLS console commands.

At the SIG business meeting Ray Carpenter (European RT-11 Sig Chairman) announced that the Mini-tasker was to become the official newsletter for the European area. Yours truly was appointed as European Correspondent to the Mini-tasker. European users can send their input to me to forward to Ken Demers. Hopefully, this wider distribution in Europe may mean that we will receive it faster - it usually takes 4 to 6 weeks.

DECUS Migration Working Group

A DECUS Working Group was formed at the symposium to examine the issues of migration, portability and compatibility. The group presently has RSX, RT and VAX representatives. The group plans to produce a roadmap of Digital operating systems, devices and options plus suggested techniques for improving program portability. Users who have ideas in this area should contact Ian Hammond, the group chairman, at the address above. We are looking for users who have had experience moving between systems.

Hidden RT-11

I have had some requests were for more information on 'hidden' RT-11 features. Here a summary of some features; perhaps other users can send additional information in this area to the Mini-tasker.

MAKE, MUNG, TECO - These three 'Standard TECO' commands are available in the RT-11/FB V4 monitor. Size constraints did not permit them to be supported in the RT-11/SJ monitor. However, a conditional change in KMON will extend support to SJ with a SYSGEN session (has anyone done this?).

PASCAL - The RT-11/FB monitor supports a PASCAL command. This command simply calls PASCAL with a copy of the command string - it is upto the PASCAL compiler to define and decode any parameters. One possible reason for a Digital RT-11 PASCAL command could be a Digital RT-11 PASCAL implementation?

RENAME - The RENAME command copies the device specification from the 'From' file to the 'To' file. That is:

RENAME DX1:A.A B.B implies DX1:(A.A B.B) implies DX1:A.A DX1:B.B

This also applies to PIP/Rename operations. The Digital comment to this was 'Yes, we finally noticed that rename operations only take place on a single device'.

COPY/FILE/DEVICE - This feature is 'buried' rather than 'hidden' - it is often overlooked. COPY/FILES/DEVICE permits volume/file copy operations which are very useful for backups.

DOLLAR - The RT-11 Keyboard Monitor will ignore a dollar sign (\$) in column zero of a console command or command file command. Its not worth reading to much into this, other than to note that it IS VAX/VMS compatible - i.e. VAX/VMS command procedures require a dollar in column zero. On the subject of command procedures, it appears that RT-11 Autopatch will contain a functional (but officially 'unsupported') version of RSX IND which handles conditionals and looping in command files. On the subject of Autopatch I guess we wont see it until RT-11 4.1 since it will probably require KMON modifications.

DIR/DUMP - RT-11 releases sometimes show evidence of work in progress. KMON has a entry for a DIR/DUMP command which is not implemented in DIR. V3 RT-11 had unused PRINT/DIR and TYPE/DIR command entries.

F4P - Since V3 there have been a couple of FORTRAN IV-plus commands in KMON. The SET FORTRAN F4P and FORTRAN/F4P commands would seem to indicate that F4P support was planned for RT-11. I guess it WAS but whether it still IS (after three years) is unknown. F4P needs an RSX-11-like RMS (Record Management System) which RT-11 doesn't have. A lot of users

feel that F4P could be (and should be) the basic compatibility tool for moving programs around Digital systems - FORTRAN-IV-minus doesnt seem to fill that position. Questions to Digital on this subject are usually answered with "ask the FORTRAN development group".

Tricky RT-11

Another category of hidden RT-11 features are distinguished by the fact that they are potentially non-portable - sometimes they are plain dirty. Portability includes moving programs between different versions of RT-11 (V3, V4, V5 etc.), different configurations (SJ, FB, SYSGEN monitors) and to and from RT-11 emulators (RSTS, planned RSX, etc.). This raises the rather fuzzy issue of the 'RT-11 Architecture' (would the people in the back rows please stop laughing!).

The most clearly defined area of the RT-11 architecture is the program request set. The Keyboard Monitor command definitions are vague - but this is not so critical. Non-EMT monitor services are precarious. But, the most clearly dangerous area are the RT-11 monitor offsets and side-affects. Offsets tend to be change between releases, configurations and implementations. Further, if a feature is not 'defined' by the RT-11 documentation then the RT development group are free to change or remove it; Murphie's law tells us that they inevitably will.

RT-11 V3 and V4 have both included some mechanisms that clean up the architecture. V3 introduced .GVAL as a portable way to get the value of a monitor offset. V3 re-implemented .DATE as an EMT for the same reason. V4 introduced the .SDTM EMT to set up the date and time - which means that a program is no longer dependent on the location or implementation of the psuedo-clock.

On the other side of the coin, RT-11 has become very dependent on job and monitor impure area offsets. V4 PIP, DIR, DUP and RESORC do a lot of bit-bashing in the monitor tables - more than V3. Now, if RT utilities need these services then its a safe bet that some user programs will also require them. But, the RT group are in a much safer position than users when it comes to betting on portability prospects.

There is a potential solution to this problem. A library of subroutines could be supplied that would handle the dirty operation on behalf of the program. Moving between implementations would just require linking to a local library. Obviously, it would be wonderful if the RT-11 group would provide such a library. If not, then perhaps we users should think about doing it.

Finally, here are two examples of RT-11 'dirty tricks':

KED - KED uses a very sneaky trick to get some scratch disk space. KED uses the final blocks of the tentative file allocated for output. Then, before it closes the file it fudges the 'highest block written' entry in the output files CSW area to ensure the closed file has the correct length. An ingenious trick - but no marks for portability. KED could have just as easily opened a temporary file.

USR COLLECT CALL - The RENAME/PROTECT and COPY operation that retains an file creation date are both implemented with a call back to a program from the USR with a register pointing at the directory entry for a file. The program (usually PIP) reads/modifies the directory entry's flag word (for PROTECT) or date word. This approach has the advantage that it doesn't cost much code in the already overcrowded USR. The USR is told that the collect call is required by setting an address in the request area odd and sending an extra argument that defines the address of the collect call subroutine. If you set that address odd by mistake you will get a crash that will be very difficult to debug.

RT-11/PDP-11 ASSEMBLY LANGUAGE TUTORIAL:
THE PHILOSOPHY AND STRUCTURE OF THE UNIBUS
Part 1 -- An Overview
by
Ray Strackbein
Cort Data Systems
77-920 Avenue of the States
Palm Desert, California 92260

THE ARCHITECTURE OF THE PDP-11

The PDP-11 family of mini-computers was originally designed to use only one bus -- the central processor unit (CPU), the memory, and all peripherals are directly attached to this bus. The primary reason for this was to allow all of the components of the computer system to 'talk' directly to each other -- not only can the CPU control its memory and all of its peripherals (as is universally the case in all computers I know of), but the all of the peripherals can 'talk' to the computer's memory and to each other as well. While it is not usually the case that one would want to have a magnetic disc talk directly to the console terminal, to do so is not prohibited by the design of the PDP-11.

The bus in a PDP-11 is called the Unibus. The Unibus transmits both address and data information (as well as other status and arbitration signals) along its length. To better understand the Unibus, let us look at just how the bus is used.

The Unibus is used to transmit information from one place in the computer system to another. Just as an example, let us take the case of memory transferring data to a disc. In order for this transfer to take place, three things must be known: 1) Where the data is coming from, 2) where the data is going to, and 3) what the data is. The data is usually a byte (8 bits) or a word (16 bits) of information and is carried along the 'DATA' lines of the Unibus in parallel form (one wire carries one bit of information -- there are 16 data wires in the Unibus: D0, D1, D2, ..., D13, D14, D15). Simple. To look at the data on the Unibus at any given time, we just look at the 16 Unibus Data wires.

Figuring out where the data is coming from and where it is going to is a bit more complex. In all data transactions there is always a bus master and a bus slave. Data can be transferred from the slave to the master, or from the master to the slave, so we must have two more wires: Data out (DATO) and Data in (DATI). If the data is out of the master the data out wire is asserted (DATO); if the data is out of the slave and into the master, the the Data in (DATI) wire is asserted. The secret to understanding this is to look at the master device. If the DATO is asserted, the the data in goes out of the master to the slave -- if the DATI (Data In) wire is asserted, then the data is going into the master from the slave.

We now set to try to figure out which device is the master and which device is the slave. The slave is easy ... its address is on the address lines of the Unibus. And the master? Well, its the device controlling the Unibus at the moment.

Now back to that example of data being transferred between disc and memory. Normally the central processor (CPU) is the bus master and memory is the slave. The CPU fetches an instruction from memory (the CPU is the master, the DATI line is

asserted to show that the data is going from the slave to the master, and the memory address is on the Unibus address lines which tells what memory location to put its data on the Unibus). If that instruction happens to enable the disc controller, the controller waits until the CPU is between bus cycles and tries to assert the bus. The bus is so constructed that if several devices try to become bus master at the same time, the device of the highest priority will become bus master. In this case, our disc drive becomes the bus master. If this is a disc write, the address of the location in memory to send the data is put on the address lines and the DATI (and MSYN) is asserted. The slave (the asserted address in memory) responds by putting the contents of its memory location in the data lines and asserts SSYN which causes the disc controller to accept the data.

All of the devices on the Unibus must have its own address on the Unibus. There are 2 to the 15th power (32,768) words which may be addressed by the Unibus (by using memory management, four times this number (131,072)). To avoid contention between memory and peripheral devices for address space, DEC has specified that the top 4K words addressable by a computer system shall be reserved for use by peripherals. Memory may use the remaining address space.

The whole result of this scheme is that there is virtually no difference between memory and peripheral devices on the Unibus. The CPU can read and write peripheral registers just as it can read and write memory. This makes things very simple both for trouble-shooting and for designing peripherals.

DESIGNING PERIPHERALS

Let's say that we want to design a disc controller to interface to the Unibus. To transfer a block of data in memory through a disc we need to know: 1) The starting memory address of the transfer, 2) The number of words to be transferred, 3) the cylinder, track, sector and drive number for the start of the transfer, 4) possibly a word to buffer the data from or to the drive, and finally, 5) miscellaneous control, status, and error information.

It is a simple matter, then, to reserve five or six 16-bit addresses for our controller to use. One word might be used for the starting memory address, another to be used as a word counter, a third for use to pass the drive, cylinder, track, and sector number between the processor and the disc controller, a fourth as a data buffer, a fifth to pass control information (to tell the drive whether to read, write, format, or seek, and whether or not to interrupt the processor when a command has completed), and a sixth for the controller to pass error and status information back to the computer.

A look through DEC's Peripherals Handbook at some of the discs will show this idea in much more detail.

MEMORY MANAGEMENT

At this point, let's deviate slightly from the topic of the UNIBUS and try to explain just a little about memory management.

You should already know that at the heart of the operation of the PDP-11 computer is general register 7, also referred to as the program counter. This register always points to the address in memory of the next computer instruction (or part of an instruction) to be executed. Since the data path of the PDP-11 is only 16 bits wide, any register (including the program counter) can only point to a 16-bit address. If

we are to try to address more than 16 bits worth of memory, some place has to be found to store some extra bits. That place is the memory management unit.

Because memory is most often addressed sequentially (especially in the case of memory which contains computer instructions) we can conclude that if the last memory reference was to location 'A', more than likely the next memory location to be addressed will be 'A+2'.

By partitioning all of the computer's memory into 32K-word partitions, we can tell the memory management unit what partition we are in and then let all references to memory refer to a relative location within that partition. Because of the large size of the 32K-word partition and because of the fact that most often memory is addressed sequentially, we will have to spend very little time juggling memory management pointers because we will usually be referring to the same partition we were referring to in the previous cycle.

This whole idea tends to fall apart whenever an interrupt occurs. There is no correlation whatsoever between the address (or partition) we were accessing before the interrupt occurred and the address we will access as a result of that interrupt. But the solution is very easy.

You should know that an interrupt causes the contents of R7 (the program counter) to be pushed on the stack. In addition, the Processor Status Word is also put on the stack by an interrupt. The PSW and R7 are then loaded with the contents of the two-word vector pointed to by that interrupt.

Because the PSW had a few extra bits which weren't being used, someone thought of using a few of those extra bits to control the memory management unit. The result is that whenever an interrupt is taken, the memory management unit is immediately notified by the bits being changed in the processor status word. This allows the processor to rapidly cross partition bounds if necessary whenever an interrupt occurs.

Because a trap instruction (BPT, IOT, EMT, and TRAP) is nothing more than an interrupt generated by the CPU, the RT-11 operating system and its layered products use the various trap instructions for calling system routines -- the memory management unit is adjusted by its hardware to account for the occurrence of an interrupt and no extensive software adjustments are necessary.

Because a DMA data transfer from a peripheral would be severely slowed if the transfer had to tell the memory management unit it wanted to stuff data in a partition different from the partition currently being used by the CPU, a peripheral can directly address 128K-words of memory through the Unibus through the memory extension bits. A DMA transfer does not use the memory management unit.

RT-11/PDP-11 ASSEMBLY LANGUAGE TUTORIAL
FALL DECUS, SAN DIEGO, NOVEMBER 7, 1980
BY
RAY STRACKBEIN
CDRT DATA SYSTEMS
77-920 AVENUE OF THE STATES
PALM DESERT, CA. 92260

```

;THESE ARE TEST PROGRAMS FOR USE IN TESTING AND DEBUGGING
;YOUR SYSTEM
;
;THIS PROGRAM WILL CAUSE A TERMINAL OR OTHER SERIAL DEVICE
;TO ECHO
;
105737 ECHO: TSTB @#177560
177560
100375 BPL ECHO
105737 1$: TSTB @#177564
177564
100375 BPL 1$
113737 MOVB @#177562,@#177566
177562
177566 BR ECHO
000766

;THIS PROGRAM WILL TEST MEMORY
;
010702 MEMTST: MOV PC,R2
005000 1$: CLR R0
011001 2$: MOV (R0),R1
005110 COM (R0)
005101 COM R1
021001 CMP (R0),R1
001007 BNE OOP
005110 COM (R0)
005101 COM R1
022001 CMP (R0)+,R1
001003 BNE OOP
020002 CMP R0,R2
103765 BLO 2$
000763 BR 1$
000000 OOP: HALT

;
;IF NOTHING RUNS, SOME PROCESSORS ALLOW INSTRUCTIONS TO BE
;STORED IN THE COMPUTER'S GENERAL REGISTERS. ALL MEMORY AND
;PERIPHERALS CAN BE REMOVED FROM THE COMPUTER AND THE COMPUTER
;CAN STILL RUN A TRIVIAL PROGRAM SUCH AS:
; ,=177700 ; THIS IS THE ADDRESS OF R0
;
000240 NOP
000777 BR 1$
1$: BR 1$
; IN THIS CASE A 240 (THE NO-OP INSTRUCTION IS PLACED IN R0,
; AND A 777 (NOTE NOT A 774) IS PLACED IN R1. IF THE ADDRESS
; OF R0 (177700) IS LOADED INTO THE SWITCH REGISTER AND
; STARTED, THE COMPUTER WILL ACTUALLY RUN WITHOUT MEMORY.
;
;
; ANOTHER USEFUL TOGGLE PROGRAM USING ONLY THE GENERAL
; REGISTERS IS A PROGRAM TO LOAD EACH MEMORY LOCATION WITH
; ITS OWN ADDRESS:
; ,=177700
005005 START: CLR R5
000240 NOP
010525 MOV R5,(R5)+
000777 1$: BR 1$
;
; IF ONE FEELS THAT THE COMPUTER'S MAIN MEMORY MIGHT BE A
; BIT UNSTABLE, ONE CAN LOAD MEMORY WITH ITS OWN ADDRESS AND
; LOOP THROUGH A READ AND COMPARE PROGRAM STORED IN MAIN
; MEMORY SUCH AS:
; ,=177700
010705 MTEST2: MOV PC,R5
005000 CLR R0
010020 MOV R0,(R0)+
020005 1$: CMP R0,R5
103775 BLO 1$
005000 2$: CLR R0
020020 3$: CMP R0,(R0)+
001003 BNE OOPS
020005 CMP R0,R5
103774 BLO 3$
000772 BR 2$
000000 OOPS: HALT

```

DEC INPUT

RT-11 TRAINING PROGRAM (Continued) CHICAGO (312) 640-5520

Canadian Customers: All registrations for U.S. courses should be made through the Ottawa (Kanata) Registrar, (613) 592-5111.

Course #	Title	Length	Start Day	tuition	Training Credit	October	November	December	January	February	March
						6 13 20 27	3 10 17 24	1 8 15 22 29	5 12 19 26	2 9 16 23 30	
EY-J2114-AO	Introduction to Minicomputers A/V	~5		\$ 400	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2016-AO	Introduction to Minicomputers	5	MO	520	1						
EY-J2116-AO	Introduction to the PDP-11 A/V	~5		450	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2024-AO	PDP-11 Assembly Language Programming	5	MO	580	1	*	*	*	*	*	*
EY-J2050-AO	Programming in MACRO-11	5	MO	630	1	*	*	*	*	*	*
EY-J2048-AO	PDP-11 Concepts	3	MO	460	1	*	*	*	*	*	*
EY-J2000-AO	RT-11 User	5	MO	630	1						
EY-J0016-AO	Programming in FORTRAN IV	5	MO	630	1		*	*	*	*	*
EY-J2002-AO	RT-11 Programmer	5	MO	630	1		*	*	*	*	*
EY-J2004-AO	RT-11 System Programmer	3	MO	520	1		*	*	*	*	*

RT-11 CUSTOMER TRAINING PROGRAM SAN FRANCISCO (408) 727-0200 Ext. 2142

Course #	Title	Length	Start Day	tuition	Training Credit	October	November	December	January	February	March
						6 13 20 27	3 10 17 24	1 8 15 22 29	5 12 19 26	2 9 16 23 30	
EY-J2114-AO	Introduction to Minicomputers A/V	~5		\$ 400	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2016-AO	Introduction to Minicomputers	5	MO	520	1		*	*	*	*	*
EY-J2116-AO	Introduction to the PDP-11 A/V	~5		450	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2024-AO	PDP-11 Assembly Language Programming	5	MO	580	1	*	*	*	*	*	*
EY-J2050-AO	Programming in MACRO-11	5	MO	630	1	*	*	*	*	*	*
EY-J2048-AO	PDP-11 Concepts	3	MO	460	1	*	*	*	*	*	*
EY-J2000-AO	RT-11 User	5	MO	630	1						
EY-J0016-AO	Programming in FORTRAN IV	5	MO	630	1		*	*	*	*	*
EY-J2002-AO	RT-11 Programmer	5	MO	630	1		*	*	*	*	*
EY-J2004-AO	RT-11 System Programmer	3	MO	520	1		*	*	*	*	*

WASHINGTON, D.C. (301) 459-7900 Ext. 2580, 2582

Course #	Title	Length	Start Day	tuition	Training Credit	October	November	December	January	February	March
						6 13 20 27	3 10 17 24	1 8 15 22 29	5 12 19 26	2 9 16 23 30	
EY-J2114-AO	Introduction to Minicomputers A/V	~5		\$ 400	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2016-AO	Introduction to Minicomputers	5	MO	520	1		*	*	*	*	*
EY-J2116-AO	Introduction to the PDP-11 A/V	~5		450	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2024-AO	PDP-11 Assembly Language Programming	5	MO	580	1	*	*	*	*	*	*
EY-J2050-AO	Programming in MACRO-11	5	MO	630	1	*	*	*	*	*	*
EY-J2048-AO	PDP-11 Concepts	3	MO	460	1	*	*	*	*	*	*
EY-J2000-AO	RT-11 User	5	MO	630	1						
EY-J0016-AO	Programming in FORTRAN IV	5	MO	630	1		*	*	*	*	*
EY-J2002-AO	RT-11 Programmer	5	MO	630	1		*	*	*	*	*
EY-J2004-AO	RT-11 System Programmer	3	MO	520	1		*	*	*	*	*

NEW YORK (212) 971-3545

Course #	Title	Length	Start Day	tuition	Training Credit	October	November	December	January	February	March
						6 13 20 27	3 10 17 24	1 8 15 22 29	5 12 19 26	2 9 16 23 30	
EY-J2114-AO	Introduction to Minicomputers A/V	~5		\$ 400	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2016-AO	Introduction to Minicomputers	5	MO	520	1		*	*	*	*	*
EY-J2116-AO	Introduction to the PDP-11 A/V	~5		450	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2024-AO	PDP-11 Assembly Language Programming	5	MO	580	1	*	*	*	*	*	*
EY-J2050-AO	Programming in MACRO-11	5	MO	630	1	*	*	*	*	*	*
EY-J2048-AO	PDP-11 Concepts	3	MO	460	1	*	*	*	*	*	*
EY-J2000-AO	RT-11 User	5	MO	630	1						
EY-J0016-AO	Programming in FORTRAN IV	5	MO	630	1		*	*	*	*	*
EY-J2002-AO	RT-11 Programmer	5	MO	630	1		*	*	*	*	*
EY-J2004-AO	RT-11 System Programmer	3	MO	520	1		*	*	*	*	*

BOSTON (617) 275-5000 Ext. 2380

Course #	Title	Length	Start Day	tuition	Training Credit	October	November	December	January	February	March
						6 13 20 27	3 10 17 24	1 8 15 22 29	5 12 19 26	2 9 16 23 30	
EY-J2114-AO	Introduction to Minicomputers A/V	~5		\$ 400	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2016-AO	Introduction to Minicomputers	5	MO	520	1		*	*	*	*	*
EY-J2116-AO	Introduction to the PDP-11 A/V	~5		450	1			CONTINUOUSLY AVAILABLE AT FACILITY			
EY-J2024-AO	PDP-11 Assembly Language Programming	5	MO	580	1	*	*	*	*	*	*
EY-J2050-AO	Programming in MACRO-11	5	MO	630	1	*	*	*	*	*	*
EY-J2048-AO	PDP-11 Concepts	3	MO	460	1	*	*	*	*	*	*
EY-J2000-AO	RT-11 User	5	MO	630	1						
EY-J0016-AO	Programming in FORTRAN IV	5	MO	630	1		*	*	*	*	*
EY-J2002-AO	RT-11 Programmer	5	MO	630	1		*	*	*	*	*
EY-J2004-AO	RT-11 System Programmer	3	WE	520	1		*	*	*	*	*

USER REQUESTS

MAX-PLANCK-INSTITUT FÜR BIOCHEMIE Ludwig Voet

8033 MARTINSRIED BEI MÜNCHEN
TELEFON (089) 8585-1
DURCHWAHL 85 85-.....
TELEX 521740 mpla d
den 24. Okt. 1980

Dear Mr. Demers,

we are collecting data under RT-11 on Mag-Tape with the MT-Handler. Now we try to read these data on VAX-11/780, but there is no direct way to do this in RT-11 or VMS Operating System via Mag-Tape. Only RX-01 Floppy-Disc is supported by FILEX under RT-11 and VMS via DOS. Due to the large data set, this is a very uncomfortable way to do this. We hope you or one of your readers know about existing written software to handle this problem.

Sincerely

L. Voet

Ludwig Voet
Max-Planck-Institut für Biochemie
D-8033 Martinsried bei München
West-Germany



1130 Hampton Avenue / St. Louis, Missouri 63139 / (314) 647-3800

Gentlemen:

I am searching for a Shop Floor Control System for a small custom cabinet shop that will run on a 64K PDT 11 under RT11, Version 4.

Sincerely,

BANK BUILDING CORPORATION

Sam Patterson
Sam Patterson
Director, Data Processing

USER RESPONSES

UNIVERSITY OF KENTUCKY
COLLEGE OF ENGINEERING
LEXINGTON, KENTUCKY 40506

WENNER-GREN
RESEARCH LABORATORY

Dear Sir:

Arun A. Ahuja, whose letter appears in Volume 6, Number 4 of the MINI-TASKER, may be interested to know about some DEC software which has been obsolete (as far as I know) since RT-11 Version 3. I refer to the LV11 support in FORTRAN/RT-11, which permits hard copy from a VT11 Display Processor. This software is described in the DEC FORTRAN/RT-11 Extensions Manual, DEC-11-LRTEA-R-D, July 1975. The manual may be available from DEC's Document Archives. I have been unable to get this software for our system because it is too old to have an SPD number. My understanding is that it is no longer available from DEC. The support library is called LVL1B.OBJ; the concatenated object modules are LVOBJ.OBJ; and the individual object files are LVA.OBJ thru LVK.OBJ, and LVU.OBJ.

Sincerely,

E. M. Lowery
E.M. Lowery
606/258-8885

Spring 81 DECUS/US Symposium
RT-11 SIG Tape

HELP WANTED

For the next symposium in Miami I need some help for building the master RT-11 tape. I am looking for a local shop that would be willing to permit me to use their equipment for building our tape. I'll need a minimum system having RX01's or RX02's and 9-track mag tape (800 bpi, 2400 ft spools). A larger disk such as an RK05, RL01 or RL02 would be very desirable. If necessary I can bring my own bootable RK05 disk cartridge. Actual use of the system would be during the evening hours. If anyone is willing to volunteer their shop and services please contact me.

Nick Bourgeois / 1738
Sandia National Laboratories
P. O. Box 5800
Albuquerque, NM 87185
(505) 844-8088

PAST SYMPOSIUM INFORMATION

Fall 80 DECUS/US Symposium
RT-11 Can Do It Workshop

This time we handled the workshop session in a more lively fashion than in the past. The workshop was run by three co-chairmen: Greg "SET SY CRASH" Adams, Ray "Gons" Strackbein and Nick "Reverend" Bourgeois. User input kept this fun forum alive and kicking for about five hours. Most of the session was taped. We'll try to have a summary for the next issue of the "Mini-Tasker".

STATUS OF RUNOFF (M02.2) FOR RT-11

This is an update on the RUNOFF M02.1 for RT-11 which I submitted to the Spring 1980 RT-11 SIG tape. I wrote in to Ken Demers shortly thereafter to follow up on that submission, and the letter was published in the June, 1980 Mini-Tasker.

First and foremost, this is a complete distribution of the latest DECUS library RUNOFF, version M02 (DECUS 11-241) as re-released and announced in DECUScore March 1980. Many, many bugs have been fixed in this version of RUNOFF, from the old M01 version. It represents Mark Lewis' (the 'M' in 'M02') last shot at it before he left for IBM land.

Besides the files needed to support RSX-11M/M+, IAS, and RSTS/E, I have added those needed to support RT-11. I have made no functional changes to RUNOFF, nor do I feel I should, as this makes the first time everybody in PDP/VAX land has had a common version of RUNOFF. I plan to submit this to the DECUS library shortly, with the fixes outlined below, as M02.2. It will also be on the Fall '80 RT-11 SIG tape.

After putting the M02.1 kit on the Spring '80 RT tape, I received a number of inquiries and a few bug reports. The main bug was that under certain conditions, RUNOFF ate up all of the stack and died. This was an omission on my part in RNRT11.MAC, where I failed to initialize the variable 'NLPG'. Also reported was the malfunctioning of 'HEADER LEVEL' when 'HEADERLEVEL' and 'HL' worked. This was fixed in the command tables, where there were some entries out of sequence. Finally, the leading and trailing formfeeds were eliminated, as this is not normal practice on RT-11, where the LP!/LS! handlers usually are set to issue their own.

There is (rumored) a DEC Standard RUNOFF coming. I don't know when they're going to announce it, but it is at least running in prototype form on a VAX at Tewksbury. I don't know how different it will be from this version.

Finally, my thanks to Chester and Lois Wilson for their nice letter and offering of the 'Australian RUNOFF', an M01 RT-11 conversion. They noted several bugs, most of which were among those fixed in the M02 release. They also outlined some very good suggestions for enhancements to RUNOFF. I will present these at the next meeting of the text processing SIG.

If anyone wishes to contact me, they may do so at:

Robert B. Denny, dba *RBD*
Creative System Design
3452 E. Foothill Blvd. Suite 601
Pasadena, CA 91107
(213) 792-9474

RT-11 V4 WITH NO DEC COMPATIBLE DISKS

By Bob Denny

I have a client who has a system with no DEC compatible disk units, but he does have a TM-11 emulating master tape unit. We got his first RT-11 system going a few years back by writing a driver for his nonstandard disk system, modifying the BSTRAP.MAC component of the RT-11 V3B monitor to include the bootstrap driver for his disk, physically taking his disk system to a site that had RK-05's, plugging it into the UNIBUS and setting a bootable system going. He takes his BPUS on master tape, and there has been no problem re-sussening since then.

Version 4 introduced some good news and some bad news. The good news is that there is no longer any device dependent code in the monitor. No more mods to BSTRAP.MAC! The bad news is that there is no way to set the distributed V4 system running unless you have a DEC compatible disk system, regardless of the distribution medium, even if you have the master tape distribution kit. The V4 DUP program won't run under earlier versions, even if you fake the version number in the monitor's offset area (it starts up but acts wierd). So there's no way to do a BOOT or COPY/BOOT of the new system, with a newly converted non-standard handler.

Not wanting to repeat the hassle of hauling his disk to the other site, we looked for a way to set going. We knew that we had to change the handler to use the deluxe macros furnished with V4 (thanks Anton!), and add the bootstrap support code. That was not too tough, and most of the bootstrap code was the same as that which we edited into V3B's BSTRAP.MAC.

Taking our best shot at that (we couldn't test it), we dove into the Software Support Manual (REALLY thanks!!!) to find out how the new V4 device independent boot system worked. It took several passes through section 7.10 to get it crystal clear. Now all we had to do was fake a 'BOOT ddn:', 'BOOT ddn:filnam', or a 'COPY/BOOT xxn:filnam ddn:' process once, and we had our running V4 distributed monitor (assuming the handler worked).

We decided that faking the COPY/BOOT would be the easiest to check out, since we could inspect the disk in blocks 0-5 after running our 'faker' routine and verify that things looked per Fig. 7-8 in the SSM. Any of the BOOT processes would just die with few clues as to the cause, if there was a problem anywhere.

Following this writeup is a listing of the V4COPYB.MAC program which was written for this purpose. We copied the basic V4 distributed monitor components and handlers, including our new system device handler, to drive 0 of the nonstandard disk, while running V3B on drive 1. Then we ran V4COPYB. Next we checked

the disk on drive 0 with DUMP to see that it looked like it should per the SSM, sec. 7.10.3.2. It did. We then hardware booted drive zero, and ... it died. After fixing a dumb error in the V4 converted driver, we replaced it on drive zero, reran V4CPYB, and hardware booted drive 0 again. This time we were greeted by the distributed SJ monitor! What fun!

If you want to know more, or have a question, feel free to write or call me. This program will be on the Fall 1980 RT-11 SIG tape.

Robert B. Denny, dba *RBD*
Creative System Design
3452 E. Foothill Blvd. Suite 601
Pasadena, CA 91107
(213) 792-9474

```
.TITLE V4CPYB
.IDENT 'Y1.0'
.ENABL LC

; This program does a COPY/BOOT, Version 4 style, while you are
; running under an earlier version, such as V3B. It is used only to
; set a first V4 system booted on a system where there are no DEC
; compatible random access mass storage devices. I used it to set
; V4 going on a system which had only compatible 9 track tape.
; NOTE:
; There are several places that must be edited prior to
; running this program. They are marked with *** style
; comments and explanation.
;
.mcall .lookup .enter .readw .writw .exit
.mcall .print .ttyin .close .hreset .fetch

; Absolute section for boot code.
;
.sect ; <<< reserve space for secondary boot >>>
. = 1000

secbot: .blkw <5*400> ; Secondary boot goes here.
        .blkw <5*400> ; First 5 blocks of monitor
pribot: .blkw <3*400> ; Space for primary boot

; ENTRY POINT.
;
.psect .pros. ; Start the relocatable code
START::
.print $stmss ; Ask user if he's sure
.ttyin ; Get Y/N
bic $177640,r0 ; Make sure it's upper case
cmp r0,$'N' ; If 'N'
beq 5$ ; exit
.fetch $hspace,$diskblk ; Fetch the handler
bcc 2$ ; If error
.print $dnfmss ; inform user
.exit ; and exit
2$: .lookup $dph,$0,$monblk ; Else, open the monitor file
bcc 10$ ; If error
.print $mnmfms ; inform user
5$: .exit ; and exit
10$: ; Else, read blocks 0-4 of monitor
```

27.

```
.readw $dph,$0,$secbot,$<5*400>,$0
bcc 20$ ; If error
.print $mremss ; inform user
.exit ; and exit
20$: ; Else, close monitor
.close $0
.lookup $dph,$0,$devblk ; Lookup the handler file
bcc 30$ ; If error
.print $hnfmss ; inform user

.exit ; and exit
30$: ; Else read primary boot
mov $pbaddr,r1 ; Copy primary boot address
asr r1 ; convert to word offset in file
bic $377,r1 ; and thence to block # where pri boot starts
swab r1
.readw $dph,$0,$pribot,$<2*400>,r1 ; Read enough to set primary boot
bcc 40$ ; If error
.print $hremss ; inform user
.exit ; and exit
40$: ; Else, move pri. boot into bes of sec. boot
.close $0 ; Close the handler file
mov $pbaddr,r1 ; Convert primary boot address to offset in
bic $177000,r1 ; the 'pribot' buffer, then to the absolute
add $pribot,r1 ; address of the start of the primary boot.
mov $secbot,r0 ; R0 --> start of secondary boot image
mov $pssize/2,r2 ; R2 = #words in primary boot
50$: mov (r1)+(r0)+,r2 ; Copy primary boot to secondary boot image
dec r2
bne 50$
mov monblk+2,$secbot+4724 ; Set up the filename for the monitor
mov monblk+4,$secbot+4726 ; in it's special spot.

; Now write out the boot data to the disk.
;
.lookup $dph,$0,$diskblk ; Non-filestructured open on the disk
.writw $dph,$0,$secbot,$400,$0 ; Primary boot to block 0
.writw $dph,$0,$secbot+1000,$<4*400>,$2 ; Sec. boot -> blks 2-5
.close $0
.exit

; Program's data base. Edit *** sections for your non standard device
; i.e., replace all 'xxn' and 'ddn' with 'SM0' and 'dd' with 'SM' to
; COPY/BOOT on non-standard device SM0:
;
pbaddr = ???? ;*** ddBOOT address from handler load map
pssize = 1000 ;*** Size of P-section ddBOOT from the
;*** MACRO-11 listings of the handler.

nlist bex
monblk: .rad50 "xxn:RT11SYS" ;*** xxn:RT11SYS is monitor to be booted
devblk: .rad50 "ddndd SYS" ;*** ddn:dd.SYS is to be system device
diskblk: .rad50 "ddn" ;*** ddn: where boot stuff goes.
dph: .blkw 6 ; DPB area for programmed requests
stmss: .asciz "COPY/BOOT xxn:RT11SYS ddm: [Y/N]?<200>" ;*** EDIT THIS TOO
dnfmss: .asciz "V4CPYB - Device handler missing"
mnmfms: .asciz "V4CPYB - Monitor file not found"
mremss: .asciz "V4CPYB - Error reading monitor file"
hnfmss: .asciz "V4CPYB - Handler file not found"
```

28.

```
hremss: .asciz 'V4CFYB - Error readings handler file'
        .list  dex
        .even
hspace:                                ; Handler does here.
        .end    start
```

DECUS Fall 80
RT-11 SIG Tape Submissions

N. A. Bourgeois, Jr.
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Systems Engineering Division 1738
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Patches through Aug 80 as published in the
"Software Dispatch" for RT-11 V4.0 and BASIC-11 V2. See
the file PATCH.DIR for more information.

Short paper describing an RT-11 application.
Presented at the "How to Solve the Operating System
Labyrinth" session. Files named RTAPPL.(RNO,DOC).

Short paper on RT-11 interrupt service routines by
the rules and otherwise or moral and immoral. Presented
at the "RT-11 Can Do It Workshop". Files named
RTISRS.(RNO,DOC).

Short paper describing a library implementation of
FORTRAN callable programmed requests for TSX/TSX-Plus.
Presented at the "RT-11 Can Do It Workshop". See
TSXLIB.DIR for more information.

MACRO-11 program to rewind and take off line all
mag tape drives on a TM11 controller. File named
RWMT.MAC.

MACRO-11 program to force the user to enter a valid
date and time. File named SIGNON.MAC.

MACRO-11 program to prevent user from interfering
with planned sequence of operations. File named
WATCH.MAC.

RTAPPL.RNO	14	15-Oct-80	RTAPPL.DOC	18	16-Oct-80
RTISRS.RNO	17	15-Oct-80	RTISRS.DOC	23	16-Oct-80
RWMT.MAC	8	07-Jul-80	SIGNON.MAC	41	07-Jul-80
WATCH.MAC	12	07-Jul-80	README.RNO	4	24-Oct-80
README.NAB	4	24-Oct-80			

9 Files, 141 Blocks

TSXLIB.DIR	1	16-Oct-80	TSXLIB.RNO	26	13-Oct-80
TSXLIB.DOC	36	16-Oct-80	TSXLIB.COM	2	08-Oct-80
MSGCOM.MAC	9	08-Oct-80	RCDLCK.MAC	10	08-Oct-80
TSXMSC.MAC	8	10-Oct-80			

7 Files, 92 Blocks

PATCH.DIR	8	16-Oct-80	RTPAT.RNO	4	06-Aug-80
BASPAT.RNO	4	26-Aug-80	RTPAT.DOC	4	06-Aug-80
BASPAT.DOC	5	26-Aug-80	DIRPAT.COM	1	04-Aug-80
DUP.001	1	05-Aug-80	DIR.001	1	05-Aug-80
LINK.001	1	05-Aug-80	LIBR.001	1	05-Aug-80
SIPP.001	1	05-Aug-80	RT11SJ.001	1	06-Aug-80
070201.COM	1	05-Aug-80	070301.COM	1	05-Aug-80
070901.COM	1	05-Aug-80	071001.COM	1	05-Aug-80
071601.COM	1	05-Aug-80	010101.COM	1	05-Aug-80
RMONFR.001	1	05-Aug-80	BSTRAP.001	1	05-Aug-80
RMUNFB.002	1	05-Aug-80	BSTRAP.002	1	05-Aug-80
010103.COM	2	05-Aug-80	BSTRAP.003	1	05-Aug-80
KMON.001	1	05-Aug-80	RMONSJ.001	1	05-Aug-80
RMONFB.003	1	05-Aug-80	MTTINT.001	1	05-Aug-80
010104.COM	1	06-Aug-80	USR.001	1	05-Aug-80
BSTRAP.004	1	05-Aug-80	RT11BL.001	1	05-Aug-80
RT11FB.001	1	05-Aug-80	RT11FB.002	1	05-Aug-80
BSTRAP.005	1	05-Aug-80	DD.001	1	05-Aug-80
DM.001	1	05-Aug-80	TM.001	1	05-Aug-80
TJ.001	1	05-Aug-80	TS.001	1	05-Aug-80
010105.COM	2	06-Aug-80	060401.COM	1	06-Aug-80
060401.COM	1	06-Aug-80	010102.COM	2	06-Aug-80
PATA1.MAC	1	25-Aug-80	PATA2.MAC	1	25-Aug-80
350102.COM	1	25-Aug-80	350103.COM	1	25-Aug-80
PATB1.MAC	1	25-Aug-80	PATB2.MAC	1	25-Aug-80
350104.COM	1	25-Aug-80	PATC1.MAC	1	25-Aug-80
PATC2.MAC	1	25-Aug-80	350105.COM	1	25-Aug-80
PATD1.MAC	1	25-Aug-80	PATD2.MAC	1	25-Aug-80
350106.COM	1	25-Aug-80	PATE1.MAC	1	25-Aug-80
PATE2.MAC	1	25-Aug-80	350107.COM	1	25-Aug-80
PATF1.MAC	1	25-Aug-80	PATF2.MAC	1	25-Aug-80
350108.COM	1	25-Aug-80	PATG1.MAC	1	25-Aug-80
PATG2.MAC	1	25-Aug-80	350109.COM	1	25-Aug-80
PATH1.MAC	1	25-Aug-80	PATH2.MAC	1	25-Aug-80
PAT11.MAC	1	25-Aug-80	350111.COM	2	25-Aug-80
PATJ1.MAC	1	25-Aug-80	PATJ2.MAC	1	25-Aug-80
PATJ3.MAC	1	25-Aug-80	PATJ4.MAC	1	25-Aug-80
350113.COM	1	25-Aug-80	PATL1.MAC	1	25-Aug-80
PATL2.MAC	1	25-Aug-80	350114.COM	1	25-Aug-80
PATM1.MAC	1	25-Aug-80	PATM2.MAC	1	25-Aug-80
350115.COM	1	25-Aug-80	PATN1.MAC	1	25-Aug-80
PATN2.MAC	1	25-Aug-80	350117.COM	1	25-Aug-80
PAT01.MAC	1	25-Aug-80	PAT02.MAC	1	25-Aug-80
350110.COM	1	25-Aug-80	PAT12.MAC	1	25-Aug-80
350202.COM	1	27-Aug-80	SUCVT1.001	1	27-Aug-80
BASPAT.COM	2	27-Aug-80	SUCVT1.002	1	27-Aug-80
350201.COM	1	27-Aug-80	DUP.002	1	21-Aug-80
FILEX.001	1	21-Aug-80	KED.001	2	02-Sep-80
KED.002	3	26-Aug-80	KED.003	2	03-Sep-80
K52.001	2	22-Aug-80	K52.002	3	22-Aug-80
K52.003	2	22-Aug-80	LINK.002	1	21-Aug-80
LP.001	1	28-Aug-80	LS.001	1	28-Aug-80
RESORC.001	1	21-Aug-80	SRCCOM.001	1	21-Aug-80
TS.002	1	28-Aug-80	061201.COM	1	05-Sep-80

061301.COM	1	05-Sep-80	062002.COM	2	05-Sep-80
070202.COM	1	05-Sep-80	070501.COM	1	05-Sep-80
070902.COM	1	05-Sep-80	071101.COM	1	05-Sep-80
071201.COM	1	05-Sep-80	170101.COM	1	05-Sep-80
170102.COM	1	05-Sep-80	0CHAIN.MAC	1	05-Sep-80
ICHAIN.KED	2	05-Sep-80	170103.COM	1	05-Sep-80
170201.COM	1	05-Sep-80	170201.TST	1	05-Sep-80
170202.COM	1	05-Sep-80	ICHAIN.K52	2	05-Sep-80
170203.COM	1	05-Sep-80	062001.COM	2	05-Sep-80
CONHOG.MAC	2	05-Sep-80	170101.TST	1	05-Sep-80
RTPAT .COM	3	05-Sep-80	170102.TST	2	05-Sep-80
170202.TST	2	05-Sep-80			

131 Files, 173 Blocks

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INDEX IS A FORTRAN CROSS REFERENCING PROGRAM. A FORTRAN SOURCE FILE PROCESSED BY INDEX WILL BE CHECKED FOR ALL OF ITS VARIABLE NAME AND LABEL USAGE. THE RESULTS WILL THEN BE LISTED IN ALPHABETICAL ORDER, LISTING ALL THE VARIABLE NAMES AND LABELS USED IN THE PROGRAM, THE LINES ON WHICH THEY WERE USED, AND HOW THEY WERE USED. IF NEEDED, THE VARIABLES FROM SPECIFIED FORTRAN SOURCE FILES CAN BE SAVED ALONG WITH THOSE OF OTHER FORTRAN SOURCE FILES AND LATER PRINTED OUT AS A SUPER INDEX GIVEING VARIABLE NAME AND THE NAMES OF ALL THE FORTRAN SOURCE MODULES IT WAS USED IN. ALSO INCLUDED IS THE CAPABILITY TO EXCLUDE FROM THE INDEX LISTING ALL VARIABLES THAT ONLY APPEAR IN A FORTRAN SOURCE MODULE IN A COMMON BLOCK ,TYPE OR EQUIVALENCE DECLARATION BUT ARE NOT USED ELSEWARE (OR LIST ONLY THOSE IF WANTED). THE USER CAN ALSO LIST ONLY THOSE VARIABLES THAT ARE GLOBAL (DEFINED IN A COMMON BLOCK) OR THOSE THAT ARE LOCAL. A VARIATION OF THE SUPER INDEX IS AVAILABLE TO LIST (ALMOST) ALL ENTRY POINTS, THE MODULES THAT THEY ARE DEFINED IN, AND ALL THE ENTRY POINTS THAT CALL THEM AND THAT THEY CALL.

INDEX WILL HANDLE ALL SOURCE FILES WRITTEN IN FORTRAN FOUR AND FORTRAN FOUR PLUS IN ACORDANCE WITH 'PDP11 FORTRAN LANGUAGE REFERENCE MANUAL' (DEC-11-LFLRA-C-D) AND THE FORTRAN 1977 STANDARD.

INDEX CAN BE GENERATED TO RUN UNDER THE RT-11 OR RSX-11M OPERATEING SYSTEMS.

README.1ST	3	29-Oct-80	INDEX .DOC	38	29-Oct-80
INDEX .HLP	3	29-Oct-80	INDEX .COM	2	29-Oct-80
IDXTKB.CMD	2	29-Oct-80	INDEX .CMD	2	29-Oct-80
LINETP.MAC	10	29-Oct-80	INDEX .MAC	9	29-Oct-80
BUFFER.MAC	5	29-Oct-80	IOLINE.MAC	16	29-Oct-80
FLOW .MAC	33	29-Oct-80	OPNCLO.MAC	25	29-Oct-80
RSX .MAC	1	29-Oct-80	STORE .MAC	10	29-Oct-80
GET .MAC	15	29-Oct-80	IOCHR .MAC	5	29-Oct-80
IMPURE.MAC	5	29-Oct-80	RAD50 .MAC	3	29-Oct-80
OUTPUT.MAC	12	29-Oct-80	TABLE .MAC	5	29-Oct-80
EVAL .MAC	37	29-Oct-80	SUPER .MAC	16	29-Oct-80

22 Files, 257 Blocks

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TAPER is a general purpose reader of 800 b.p.i. magnetic tapes. Its specialty is to take character data in an unknown format and, through interaction with the user, write the data in Digital Equipment format to any write-enabled disc. TAPER may be used to read "foreign" tapes.

TECHNICAL INFORMATION

The user interface and formatting is handled by a program written in FORTRAN. The tape handler is written in MACRO-11. Only the FORTRAN programs, the EBCDIC collating sequence file, and documentation with the name TAPER was written by the First National Bank of Chicago. All EXTMT routines were obtained from the Digital Equipment Computer User's Society. EXTMT was written by N. A. Bourgeois of Scandia Laboratories and operated there for the United States Department of Energy.

05-Nov-80				
TAPER .DOC	14	30-Oct-80	TAPER .HLP	16
TAPER .FOR	13	30-Oct-80	TAPSUB.FOR	9
TAPER .OBJ	16	30-Oct-80	TAPSUB.OBJ	16
TAPER .SAV	52	30-Oct-80	EXTMT .DOC	3
EXTMT .MAC	23	30-Oct-80	EXTMT .OBJ	2

10 Files, 164 Blocks

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A FORTRAN IV/RT-11 Program that will renumber statement lines in a FORTRAN IV source program. (RT-11, V3B)

LOOKUP.FOR	5	23-Oct-80	LOOKUP.OBJ	10	23-Oct-80
NCLOSE.FOR	1	23-Oct-80	NCLOSE.OBJ	4	23-Oct-80
NREAD .FOR	2	23-Oct-80	NREAD .OBJ	5	23-Oct-80
NIF .FOR	2	23-Oct-80	NIF .OBJ	6	23-Oct-80
NGOTO .FOR	2	23-Oct-80	NGOTO .OBJ	6	23-Oct-80
NACCEP.FOR	2	23-Oct-80	NACCEP.OBJ	4	23-Oct-80
LARRY1.OBJ	38	15-Nov-80	RESEQ .SAV	48	15-Nov-80
RESEQ .FOR	17	15-Nov-80	RESEQ .OBJ	25	15-Nov-80

16 Files, 177 Blocks

333 Cedar Street
New Haven, CT 06150
203-436-3617

Utility routines for Fortran/RT11.

05-Nov-80
DSLIB .DOC 56 30-Oct-80 DSLIB .FOR 18 30-Oct-80
ASLOOK.MAC 6 30-Oct-80 DSLIB .COM 1 30-Oct-80
4 Files, 81 Blocks

Mr. Robert L. Sharp
7073 Jasper Drive
Middletown, MD 21769

A Polish Interpreter.

POLISH.MAC 83 02-Nov-80
1 Files, 83 Blocks

Hal Hovland
Hal Systems Corp
3544 SW 172nd
Seattle, WA 98166
(206) 244-6606

FORTTRAN callable macros to:
Move word, Move byte, ASCII/integer conversions, shift bits.

HALIB .MAC 20 03-Nov-80
1 Files, 20 Blocks

Robert Abramson
Digital Equipment Corp.
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Nashua, NH 03061
603-884-8333

Summary description of PDP-11 FORTRAN OTS

README.DAT 1 27-Oct-80 F40TS .DOC 439 27-Oct-80
2 Files, 440 Blocks

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Yale Medical School Dep't Pharmacology

Robert B. Denny
Creative System Design
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Complete distribution of DECUS Standard RUNOFF, Version M02.2
Includes RT-11 support, as well as RSX-11M, IAS and RSTS/E.
This is the latest version, with many enhancements and bug
fixes from older versions. It is the Text Processing SIG's
standard version, previously not available for RT-11. Has
RT-11 .SAV file, ready to run.

RNOASM.CMD	RNOBLD.CMD	RNOIAS.CMD	RSTASH.CMD
RSTBLD.CMD	RNOASM.COM	RNOLNK.COM	RNOBLD.CTL
RUNOFF.DOC	CMTAB .MAC	COMND .MAC	ERMMSG .MAC
FMTCM .MAC	HYPHEN.MAC	INDEX .MAC	PINDX .MAC
RNCMD .MAC	RNFIO .MAC	RNORSX.MAC	RNPRES .MAC
RNRT11.MAC	RSTS .MAC	RT11 .MAC	RUNOFF.MAC
START .MAC	UARITH.MAC	SMAC .MLS	RBLDFC.ODL
RNO .ODL	RNOBLD.ODL	RNOIAS.ODL	RUNOFF.RNO
RUNOFF.SAV			

RT-11 conversion of David Sykes' excellent RATFOR preprocessor.
This is a slightly earlier version than is in use on RSX-11M,
but it is just fine for RT-11. No RT-specific documentation
(sorry), read the comments at the head of RATRT.RAT for RT-11
usage information. Also has nice runtime library. OK for RT-11
FORTRAN V2.5 on inline code. Malfunctions on earlier FORTRAN
V2's in inline code.

	DIST .COM	RATBLD.COM	GETSTR.FOR
PUTSTR.FOR	RATRT .FOR	RAT1 .FOR	RAT2 .FOR
RAT3 .FOR	STRLIB.FOR	RATHLP.MAC	RATFOR.MAP
STRLIB.OBJ	CCHAR .RAT	CDATIM.RAT	CDEFIO.RAT
CFOR .RAT	CFUNC .RAT	CLINE .RAT	CLIST .RAT
CLOOK .RAT	COUTLN.RAT	CPRTL.N.RAT	CSTR .RAT
CUCLC .RAT	DEFIN .RAT	DEFINS.RAT	RATDEF.RAT
RATRT .RAT	RAT1 .RAT	RAT2 .RAT	RAT3 .RAT
STRLIB.RAT	TEST1 .RAT	TEST2 .RAT	TEST3 .RAT
TEST4 .RAT	TEST5 .RAT	CFILES.RT	RATFOR.SAV
RATFOR.DOC	RATFOR.RNO		

UNOFFICIAL versions of DECUS 'C' system. Submitted to whet
your appetite. The full official source distribution is available
from the Structured Languages SIG and the DECUS Library. Look at
XRF.COM for usage example. Implements most UNIX library funct's
per Kernighan and Ritchie. Put *.(SAV,OBJ) and STDIO.H on a
disk and assign it as device "C:" before use. The documentation
included is PRELIMINARY also!!! CAVEAT!!! If you're at all
interested in this, GET THE OFFICIAL TAPE.

		CC .SAV	AS .SAV
XR .SAV	SUPORT.OBJ	CLIB .OBJ	STDIO .H
XRF .H	XRFI .C	XRFD .C	XRFO .C
XRF1 .C	XRF2 .C	XRF3 .C	XRF .COM
CSYST .DOC	CC .DOC		

SUBMIT.DOC	6	02-Nov-80	RNOASM.CMD	2	02-Nov-80
RNOBLD.CMD	2	02-Nov-80	RNOIAS.CMD	3	02-Nov-80
RSTASM.CMD	1	02-Nov-80	RSTBLD.CMD	2	02-Nov-80
RNOASM.COM	1	02-Nov-80	RNOLNK.COM	1	02-Nov-80
RNOBLD.CTL	4	02-Nov-80	RUNOFF.DOC	146	02-Nov-80
CMTAB .MAC	12	02-Nov-80	COMND .MAC	8	02-Nov-80
ERMSG .MAC	10	02-Nov-80	FMTCH .MAC	26	02-Nov-80
HYPHEN.MAC	55	02-Nov-80	INDEX .MAC	11	02-Nov-80
PINDX .MAC	9	02-Nov-80	RNCMD .MAC	21	02-Nov-80
RNFIO .MAC	9	02-Nov-80	RNORSX.MAC	6	02-Nov-80
RNPRES .MAC	5	02-Nov-80	RNRT11.MAC	30	02-Nov-80
RSTS .MAC	1	02-Nov-80	RT11 .MAC	1	02-Nov-80
RUNOFF.MAC	74	02-Nov-80	START .MAC	9	02-Nov-80
UARITH.MAC	3	02-Nov-80	SMAC .MLS	6	02-Nov-80
RBLDFC.ODL	2	02-Nov-80	RNO .ODL	1	02-Nov-80
RNOBLD.ODL	2	02-Nov-80	RNOIAS.ODL	1	02-Nov-80
RUNOFF.RNO	122	02-Nov-80	RUNOFF.SAV	33	02-Nov-80
DIST .COM	3	02-Nov-80	RATBLD.COM	3	02-Nov-80
GETSTR.FOR	2	02-Nov-80	PUTSTR.FOR	2	02-Nov-80
RATRT .FOR	65	02-Nov-80	RAT1 .FOR	43	02-Nov-80
RAT2 .FOR	41	02-Nov-80	RAT3 .FOR	49	02-Nov-80
STRLIB.FOR	70	02-Nov-80	RATHLP.MAC	5	02-Nov-80
RATFOR.MAP	16	02-Nov-80	STRLIB.OBJ	95	02-Nov-80
CCHAR .RAT	2	02-Nov-80	CDATIM.RAT	1	02-Nov-80
CDEFIO.RAT	1	02-Nov-80	CFOR .RAT	1	02-Nov-80
CFUNC .RAT	1	02-Nov-80	CLINE .RAT	1	02-Nov-80
CLIST .RAT	2	02-Nov-80	CLOCK .RAT	1	02-Nov-80
COUPLN.RAT	1	02-Nov-80	CPRTLN.RAT	1	02-Nov-80
CSTR .RAT	1	02-Nov-80	CUCLC .RAT	1	02-Nov-80
DEFIN .RAT	6	02-Nov-80	DEFINS.RAT	4	02-Nov-80
RATDEF.RAT	6	02-Nov-80	RATRT .RAT	54	02-Nov-80
RAT1 .RAT	34	02-Nov-80	RAT2 .RAT	32	02-Nov-80
RAT3 .RAT	38	02-Nov-80	STRLIB.RAT	61	02-Nov-80
TEST1 .RAT	1	02-Nov-80	TEST2 .RAT	1	02-Nov-80
TEST3 .RAT	2	02-Nov-80	TEST4 .RAT	1	02-Nov-80
TEST5 .RAT	1	02-Nov-80	CFILES.RT	2	02-Nov-80
RATFOR.SAV	77	02-Nov-80	RATFOR.DOC	104	02-Nov-80
RATFOR.RNO	95	02-Nov-80	CC .SAV	97	02-Nov-80
AS .SAV	46	02-Nov-80	XR .SAV	23	02-Nov-80
SUPPORT.OBJ	1	02-Nov-80	CLIB .OBJ	76	02-Nov-80
STDIO .H	5	02-Nov-80	XRF .H	4	02-Nov-80
XRF1 .C	9	02-Nov-80	XRFD .C	2	02-Nov-80
XRF0 .C	12	02-Nov-80	XRF1 .C	12	02-Nov-80
XRF2 .C	11	02-Nov-80	XRF3 .C	7	02-Nov-80
XRF .COM	1	02-Nov-80	CSYSTM.DOC	466	02-Nov-80
CC .DOC	109	02-Nov-80			

91 Files, 2433 Blocks

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SPOOL/RT-11 V1.01, AUTOQ-automatic queuing utility.
Based on the paper "A Transparent Spooler for RT-11" given in

Chicago 80, SPOOL/RT-11 V1.01 is an updated version which includes documentation.

DSC/RT-11 is a verifying disk scanner which reports the unique and common files found on the file structured devices (disk & mas tape).

SPRELN.DOC	13	17-Oct-80	SPUSER.DOC	37	17-Oct-80
SP .MAC	11	17-Oct-80	SPV3B .MAC	13	17-Oct-80
SPOOL .MAC	9	20-Apr-80	SPINT .MAC	26	23-Aug-80
SPINIT.MAC	25	17-Oct-80	SPOOL .REL	8	17-Oct-80
SPOOLX.REL	7	17-Oct-80	SP .SYS	2	17-Oct-80
SPX .SYS	2	17-Oct-80	SPBLD .COM	1	17-Oct-80
SPBLDX.COM	1	17-Oct-80	SPBLDL.COM	1	17-Oct-80
DSC .SAV	24	01-Nov-80	DSC .RNO	14	01-Nov-80
DSC .DOC	19	01-Nov-80	DSCLIN.PAS	3	28-Aug-80
DSCASC.PAS	3	21-Aug-80	DSCCMP.PAS	4	27-Aug-80
DSCPRN.PAS	4	28-Aug-80	DSCCSI.PAS	12	01-Nov-80
DSCMAN.PAS	5	01-Nov-80	DSCHED.PAS	4	01-Nov-80
DSCSET.PAS	9	01-Nov-80	DSCCMP.COM	1	15-Sep-80
DSCTST.COM	1	25-Aug-80	DSCLNK.COM	1	27-Aug-80
DSCMAC.COM	1	31-Oct-80	README.TXT	2	01-Nov-80
AUTOQ .OBJ	3	22-Mar-80	AUTOQ .REL	11	16-Mar-80
AUTOQ .MAC	14	22-Mar-80			

33 Files, 291 Blocks

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Miscellaneous Handlers and formatters.

PH .MAC	33	05-Nov-80	LPTT40.MAC	21	05-Nov-80
HS .MAC	26	05-Nov-80	CLEAN .MAC	12	05-Nov-80
LPOUT .FOR	5	05-Nov-80	STRIP .MAC	2	05-Nov-80
PH .DOC	3	05-Nov-80	README.DOC	7	05-Nov-80
CLEAN .DOC	4	05-Nov-80	LPOUT .DOC	5	05-Nov-80
LPTT40.DOC	8	05-Nov-80	HS .DOC	16	05-Nov-80

12 Files, 142 Blocks

See the next issue of this newsletter for information as to how you may obtain this RT-11 Fall '80 Symposium Tape.



SYSCON DESIGN INC.

815 MANHATTAN AVE. Suite F
MANHATTAN BEACH, CA. 90266
(213) 372-8544

FMX - File Management System
GP-10, DL100 Graphics Software

FMX is a disk-based package of FORTRAN-callable routines which provide an ISAM (Indexed Sequential-Access Method) file capability to the FORTRAN programmer. The multilevel directory structure of FMX enables the number of disk accesses per data record retrieved to be cut to the absolute minimum while allowing very efficient allocation of disk space. Data records consist of key fields and data fields, where the data fields may be variable length. FMX supports random, sequential, and mixed access modes, and provides special routines for queue maintenance and partitioned files.

GP-10 is a package of FORTRAN-callable subroutines which provide a basic yet versatile graphics capability to the FORTRAN programmer. GP-10 will operate any graphics device accepting the Tektronix character codes, which include the Lear-Siegler ADM-3A with a Retro-Graphics board, and the Retro-Graphics VT-100. The basic MOVE, DRAW, and POINT routines are available in device and virtual (user-defined) coordinates, and in absolute and relative modes. In addition, there are utilities for displaying text, defining dashed lines, and doing interactive graphics with such features as cursor control and selective erasure.

DL-100 is an interactive graphics program which enables the user to construct and annotate graphs of data stored on a computer file using English-like commands. Data files may be entered manually or generated by programs written in any language. The user may choose from a number of plot types, define axes, define symbols to highlight plotted points, and display text horizontally and vertically.

FMX, GP-10, and DL100 are currently available for DEC PDP-11 computers using RT-11 and TSX operating systems. For more information contact SYSCON DESIGN INC., 815 Manhattan Avenue, Suite F, Manhattan Beach, CA 90266; telephone (213) 372-8544.

OPERATING SYSTEM RT11	VERSION V3,V3B,V4	SYSTEM PROGRAM OR DOCUMENT TITLE DUP.SAV	VERSION OR DOCUMENT PART NO. V3.04B and V4.00C	DATE 3 NOV 80
NAME: DOUGLAS RANDALL Division of FIRM: Walter Kidde & Co. Inc. 6 Pawcatuck, C Pawcatuck, CT 02891			DEC OFFICE Meridan, CT	DO YOU HAVE SOURCES? YES <input type="checkbox"/> NO <input checked="" type="checkbox"/>
ADDRESS:			REPORT TYPE/PRIORITY	1. <input type="checkbox"/> HEAVY SYSTEM IMPACT
CUST. NO.: 40608			<input type="checkbox"/> PROBLEM/ERROR	2. <input checked="" type="checkbox"/> MODERATE SYSTEM IMPACT
SUBMITTED BY: Earl T. Ellis Jr.			<input type="checkbox"/> SUGGESTED ENHANCEMENT	3. <input type="checkbox"/> MINOR SYSTEM IMPACT
PHONE: (203) 599-1750			<input checked="" type="checkbox"/> OTHER	4. <input type="checkbox"/> NO SIGNIFICANT IMPACT
ATTACHMENTS			5. <input type="checkbox"/> DOCUMENTATION/SUGGESTION	
MAG TAPE <input type="checkbox"/> FLOPPY DISKS <input type="checkbox"/> LISTING <input type="checkbox"/> DECTAPE <input type="checkbox"/>			CAN THE PROBLEM BE REPRODUCED AT WILL? YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>	
OTHER:			COULD THIS SPR HAVE BEEN PREVENTED BY BETTER OR MORE DOCUMENTATION? YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>	
CPU TYPE PDP11/34A	SERIAL NO. 44859	MEMORY SIZE 256Kb	DISTRIBUTION MEDIUM RL01	SYSTEM DEVICE RL01
			DO NOT PUBLISH <input type="checkbox"/>	

1. Differences in DUP.SAV from V3 (and V3B) and V4 of RT11 which are not clearly documented and can result in severe data destruction. This has to do with the /I option of DUP.

- A. Under RT-11 V3 and V3B, /I had 3 optional entries (/I [rstart:rstop:wstart]) which allowed for Block Transfer from one device to another. In my case, I had a .COM file which updated files without copying the whole file. The line looked like this:

```
R DUP (V3.04B)
*DL1:DB2.DAT-DY1:/I:16:161:150/Y
*
```

When, by mistake, this line was run under V4, the DY was copied in Image mode destroying the directory of the RL, and leaving it as a copy of the DY, and only 900+ blocks lonk!

- B. After re-reading sections of the SYSTEM RELEASE NOTES (Page 2.8) and the SYSTEM USERS MANUAL (Page 8.5) I have learned that this COM file should be edited to read something like this!

```
R DUP (V4.00C)
*DL1:/G:150-DY1:/I/G:16/E:161/Y
*
```

That is a rather major change, and I didn't even get a glimmer of this when I read the documents the first time. When one is using IMAGE COPY, Every keystroke that changes is very important.

CSIRO

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RNC:CR
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SPR Section. Fortran IV V2.1 Inline Code

In the previous issue of Mini-tasker the problem of the Inline code from the V2.1 compiler was raised by Mark Liverman. I have identified the problem and submitted an SPR. The reply I received was, however, rather common: "... the next release ..." etc. I enclose a copy of the SPR response.

Put simply, the problem of faulty code generation by the Inline section of the compiler is a deep rooted one of poor design in the register optimisation phase. The failure is quite reproducible. Given any situation where two integer variables are in use close together, and in particular if they are used as array indices, the optimiser may attempt to keep them both in registers rather than in memory. This is arguably an undesirable practice in a real-time Fortran system with user-written interrupt-driven processes going on. Anyhow, the optimiser can fail and get the variables tangled. Of course, if you jumble the statements around a bit, you may be able to fool the optimiser into not trying. I have found the following line to be an excellent example of this:

Type *, I, J, A(I,J).

With optimisation for space, this will be wrongly compiled and will produce quite erroneous output for A(I,J) at least. Examination of the generated code usually shows R0 and R1 being confused.

I am afraid that the only safe answer is to go back to using threaded code. The drawback is a small loss of speed. The advantages are reliable code, smaller core images for large programs, more error checking, a smaller compiler, etc. Alternatively, as noted in the SPR reply, try using /M:CSE with inline code. This has the effect of turning off the optimiser section concerned with common subexpressions.

Yours sincerely,



R.N. Caffin

SUBJECT: SPR NUMBER 11-28332

	SYS	PROD
SOFTWARE:	SYSTEM VERS PRODUCT	VERS COMPONENT
	RT-11 3B FORTRAN-IV	2.1 COMPILER

PROBLEM
STATEMENT

Fortran compiler optimizes incorrectly the enclosed subroutine.

RESPONSE

The problem with the compiler will be fixed in the next release. In the meantime as a temporary workaround, please compile the program using the threaded code option or use the /M:CSE option when compiling with inline code.

OPERATING SYSTEM RT-11SJ	VERSION V04	SYSTEM PROGRAM OR DOCUMENT TITLE PIP	VERSION OR DOCUMENT PART NO.	DATE 15-Aug-80
NAME: E. L. Friedman FIRM: Montreal Woollens Ltd.		DEC OFFICE	DO YOU HAVE SOURCE? YES <input type="checkbox"/> NO <input checked="" type="checkbox"/>	
ADDRESS: 291 Elgin St. N, Cambrdg, Ont		REPORT TYPE/PRIORITY	1. <input type="checkbox"/> HEAVY SYSTEM IMPACT 2. <input type="checkbox"/> MODERATE SYSTEM IMPACT 3. <input checked="" type="checkbox"/> MINOR SYSTEM IMPACT 4. <input type="checkbox"/> NO SIGNIFICANT IMPACT 5. <input type="checkbox"/> DOCUMENTATION/SUGGESTION	
CUST. NO.:	PHONE: 519-621-5460	CAN THE PROBLEM BE REPRODUCED AT WILL? YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>		
SUBMITTED BY: D. Friedman		COULD THIS SPR HAVE BEEN PREVENTED BY BETTER OR MORE DOCUMENTATION? YES <input type="checkbox"/> NO <input type="checkbox"/> PLEASE EXPLAIN IN PROVIDED SPACE BELOW.		
ATTACHMENTS MAG TAPE <input type="checkbox"/> FLOPPY DISKS <input type="checkbox"/> LISTING <input type="checkbox"/> DECTAPE <input type="checkbox"/>				
OTHER:	CPU TYPE 11/40	SERIAL NO.	MEMORY SIZE 32K	DISTRIBUTION MEDIUM RX05
			SYSTEM DEVICE RE05	DO NOT PUBLISH <input type="checkbox"/>

When using a sysgen'd monitor with multi-terminal support the command:

COPY TT: FILE does not recognize Ctrl/Z as end of file and loops.

SPR NUMBER: 11-32528

DATE: 23-OCT-80

SOFTWARE SYSTEM(S) AFFECTED: RT-11

VERSION(S): V4.0

SOFTWARE COMPONENT(S) AFFECTED: SJ MTTY MONITOR

VERSION(S): V4.0

STATEMENT: Issuing CTRL/Z to end file created by COPY TT: FILE under SJ multi-terminal monitor causes system to hang.

RESPONSE: Thank you for bringing your problem to our attention. Upon investigating your problem we found other errors in the SJ terminal service routine. First the prompt for input is not printed until after the first character is typed. A second problem is typing a CTRL/C does not abort the operation and actually corrupts the input. An article is being generated for a future issue of the Software Dispatch to correct the above errors.



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etc.